# Large-scale compression of software source code

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### Why data compression?

#### It provides rich dividends:

- Space saving, of course!
  - ... and thus, reduced storage and transmission costs
- Energy saving: less servers involved
- Speed: not so obvious!
  - better exploitation of the memory hierarchy
  - reduced burning out of sectors in SSD disks
  - no need to decompress the whole data, in some approaches





The edition of 2016

#### **Brotli: A General-Purpose Data Compressor**

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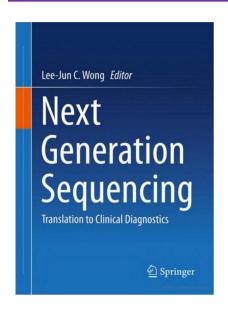
table of contents doi>10.1145/3231935



## Searching Compressed Data



We were the first in 2000 to show how to search compressed data, without decompressing all of them



Fast and accurate short read alignment with Burrows–Wheeler transform

HLi, R Durbin - bioinformatics, 2009 - academic.oup.com

Motivation: The enormous amount of short reads generated by the new DNA sequencing technologies call for the development of fast and accurate read alignment programs. A first generation of hash table-based methods has been developed, including MAQ, which is ...

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[HTML] Ultrafast and memory-efficient alignment of short DNA sequences to the human genome

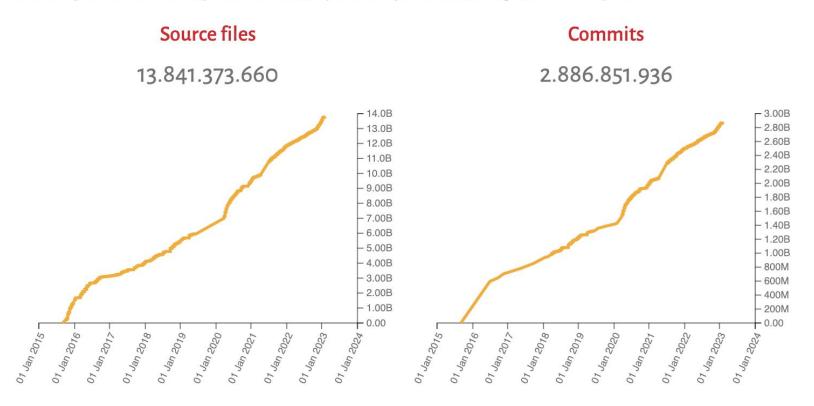
<u>B Langmead</u>, <u>C Trapnell</u>, <u>M Pop</u>... - Genome ..., 2009 - genomebiology.biomedcentral.com Bowtie is an ultrafast, memory-efficient alignment program for aligning short DNA sequence reads to large genomes. For the human genome, Burrows-Wheeler indexing allows Bowtie to align more than 25 million reads per CPU hour with a memory footprint of approximately



### The Software Heritage...

#### Size

As of today the archive already contains and keeps safe for you the following amount of objects:



Raw space (files) is about 1 PetaByte (10<sup>15</sup> bytes)

### How to approach the problem?

### Squash Compression Benchmark

The Squash library is an abstraction layer for compression algorithms, making it trivial to switch between them... or write a benchmark which tries them all, which is what you see here!

The Squash Compression Benchmark currently consists of 28 datasets, each of which is tested against 29 plugins containing 46 codecs at every compression level they offer—the number varies by codec, but there are 235 in total, yielding 6,580 different settings. The benchmark is currently run on 9 different machines for a current grand total of 59,220 configurations, and growing.

SKIP TO RESULTS (PRETTY PICTURES!)

LEARN MORE ABOUT SQUASH 🗹

# Three main families of compressors

[1977] Lempel-Ziv parsing and its derivatives

- The classic approach, made famous by gzip
- A revamped interest around it: 7z, zstd, brotli, lzfse, ...

[1994] Burrows-Wheeler Transform and its derivatives

- Better entropy-bounds than gzip-like compressors (□ bzip)
- But building a big BWT is costly and complicated...

[2000] Compressed (self-)indexes: offer space/decompress trade-offs

FM-index, CSA, LZ-trie,...

[Ferragina-Manzini, J.ACM '05]

- In the compressed space, they include text + (full-text) index
- Query/Decompression time is theoretically optimal

New compressors are based on sophisticated ML models, but slow.

### Something to be careful...

We are talking about compressing a:

collection of files of different types.

We have to apply a different compression scheme:

**PPC: Permute + Partition + Compress** 

and study it in two experimental scenarios

- BackUp for the streaming access to compressed data
- Random access to single compressed files

#### Two main issues to deal with:

- Permuting based on which information? filename, path, file type, content,...
- Partitioning based on which block size? BackUp/Random access, compressor,...

### Our experimental framework

The C and Python repositories in Github with most stars, each consisting of about 25 Gb.











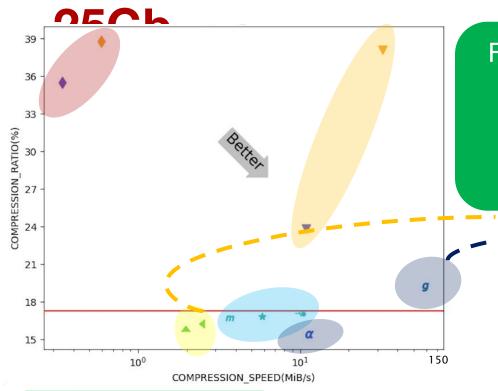






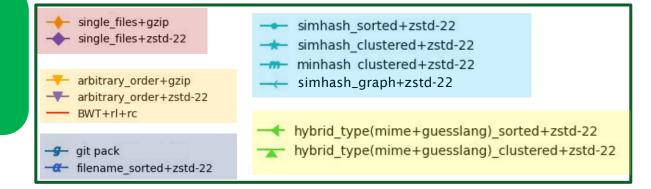


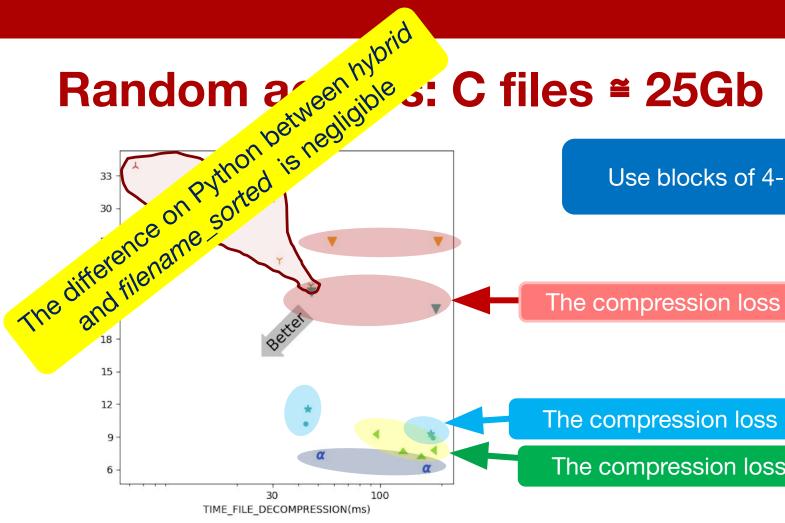
### BackUp scenario: Python files ≅



For random access, even with 4 MB blocks, HYBRID is very robust by loosing only about 1% in compression ratio wrt Backup scenario

On C files we get about 6% vs 27% compr. ratio of gzip





Use blocks of 4-16 MBs

The compression loss is about 10%

The compression loss is about 3%

The compression loss is about 1%

```
arbitrary order+blocks+gzipblock size = 4MiB
                                                                                                                                          csa_wt<wt_huff>rrr vector < 31 >
                                                                 simhash_sorted+blocks+zstd-22block size = 4MiB
      arbitrary order+blocks+gzipblock_size = 16MiB
                                                                                                                                          csa_wt<wt_huff>rrr vector < 63 >
                                                                 simhash sorted+blocks+zstd-22block size = 16MiB
                                                                 simhash clustered+blocks+zstd-22block size = 4MiB
                                                                                                                                          csa_wt<wt_huff>rrr vector < 127 >
      arbitrary order+blocks+zstd-22block size = 4MiB
                                                                 simhash_clustered+blocks+zstd-22block size = 16MiB

    csa sada

      arbitrary order+blocks+zstd-22block size = 16MiB
                                                                hybrid type(mime+guesslang) sorted+blocks+zstd-22block size=4MiB
                                                                hybrid type(mime+guesslang)_sorted+blocks+zstd-22block size=16MiB
     filename sorted+blocks+zstd-22block size = 4MiB
                                                                hybrid_type(mime+guesslang)_clustered+blocks+zstd-22block size = 4MiB
filename sorted+blocks+zstd-22block size = 16MiB
                                                                hybrid_type(mime+guesslang)_clustered+blocks+zstd-22block_size = 16MiB
```

#### The moral...

We compared **Data-aware-PPC** versus **Context-aware** approaches (à la GitPack), achieving the following results:

- Our features subsume the "context" at a larger decompression speed than GitPack
- They are robust with respect to block size
- zstd is a good compressor: large window at 100MB/sec speed

#### Next, we aim at:

- Investigating other orderings and partitionings
- □ C++ coding & HPC is needed because a single machine would take about 3 months of computation (≅ 200 Tb space at about 4k euro).

It's time to try it on the whole SI

